## Optimization and Uncertainty

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Prof. Dr. Martin Hoefer Conrad Schecker, Lisa Wilhelmi



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## Assignment 3

**Exercise 3.1** PROPHET with multiple items

For the PROPHET problem, let  $p_i$  be the probability that  $v_i$  is optimal, where i = 1, ..., n. Furthermore, suppose that  $\tau_i$  is such that  $\Pr[v_i \ge \tau_i] = p_i$ , i.e., the  $p_i^{\text{th}}$  percentile for  $v_i$ . For simplicity, assume that such a threshold  $\tau_i$  always exists. Define

$$\tilde{v}_i(p_i) := \mathbb{E}[v_i \mid v_i \ge \tau_i]$$

as the expected value of  $v_i$  given that it lies in the top  $p_i^{\text{th}}$  percentile. We consider the following strategy to pick one item: When item i shows up, if no item has been chosen among 1, ..., i-1, reject it with probability 1/2 outright, else accept it if  $v_i \geq \tau_i$ .

- a) Show that the algorithm achieves a value of at least  $\frac{1}{4} \cdot \mathbb{E}[v_{\max}]$ .
- b) Suppose we are now allowed to choose up to k out of the n items. Thus, the sequence ends with acceptance of the k-th item or after the arrival of the n-th item. The quantity of interest is given by  $\mathbb{E}[\text{sum of values of best } k \text{ items}]$ . Show that the algorithm achieves at least  $\frac{1}{4}$  of it. *Hint: Redefine*  $p_i$  as the probability that  $v_i$  is among the top k values.
- c) If the distributions are finite, then a threshold  $\tau_i$  as defined above might not exist. Can you adjust the algorithm to handle this case? Explain your answer.

Exercise 3.2 Two-round IID Prophets

Consider two rounds of an IID PROPHET problem. The finite distribution  $\mathcal{D}$  has support  $\{0, 1, n\}$ , n > 1. The probabilities are  $\Pr[v_1 = 1] = q$ ,  $\Pr[v_1 = n] = (1-q)/n$ , and  $\Pr[v_1 = 0] = (1-1/n)(1-q)$ . The same holds for  $v_2$ .

- a) Let  $ALG^*$  be the optimal online algorithm for this instance. Determine  $\mathbb{E}[v(ALG^*)]$  as a function of q and n.
- b) Derive a lower bound on the competitive ratio as a function of q and n. What are (asymptotically) optimal choices for q and n when we want to maximize this lower bound? What is the resulting bound? Prove your answer.

(3+3+1=7 points)

(3 + 4 = 7 points)

## Exercise 3.3 Markov Decision Processes

Consider a stochastic decision problem similar to the one with the envelopes solved in the lecture. There are n boxes where box i contains a prize of 1 Euro with probability  $q_i$  and is empty otherwise. The game ends when a non-empty box is opened. At each point in time, one can also decide to stop playing. This means, the final prize is either 0 Euros or 1 Euro. Opening box i costs  $c_i$  Euros. The goal is to maximize the final prize minus the costs for opening.

- a) Model this problem as a Markov decision process. In particular, give the state and action sets as well as transition probabilities and rewards.
- b) Derive an optimal policy for this problem.

The assignments and further information on the course are provided on our website: https://algo.cs.uni-frankfurt.de/lehre/oau/sommer23/oau23.shtml

Contacts: {schecker,wilhelmi}@em.uni-frankfurt.de.