

## Assignment 7

Issued: 08.06.2021  
Due: 22.06.2021, **10:00h**

The handling time of this assignment is **two weeks**. It will be discussed in the exercise session on 25.06.2021. There will be **no exercise session on 18.06.2021**.

This is the first assignment counting for the **second part** of the lecture.

### Exercise 7.1 *Direct and persuasive schemes* (3 points)

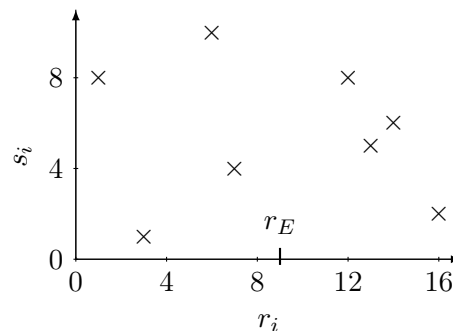
Show Proposition 3 from the lecture notes: *For every signaling scheme  $\varphi$  there is a direct and persuasive scheme  $\varphi'$  such that  $\mathcal{S}$  obtains the same expected value in  $\varphi$  and  $\varphi'$ .*

### Exercise 7.2 RANDOM-ORDER SIGNALING (3 + 4 + 4 points)

In RANDOM-ORDER SIGNALING,  $n$  prize-pairs are revealed to both the sender  $\mathcal{S}$  and the receiver  $\mathcal{R}$  at the beginning. Afterwards, the prize-pairs are packed into  $n$  boxes. The boxes are closed (which makes them look identical), permuted in uniform random order, and then labeled from 1 to  $n$ . Now  $\mathcal{S}$  can look into all boxes and then sends a signal to  $\mathcal{R}$ .

- Show that the optimal scheme  $\varphi^*$  for  $\mathcal{S}$  is symmetric.
- Describe an algorithm to compute  $\varphi^*$  in polynomial time. Give an informal argument why your algorithm yields an optimal policy indeed.  
*Hint: For starters, consider the geometric visualization of the example instance below. The instance consists of eight boxes with prize-pairs (marked by crosses) as shown in the table. As introduced in the lecture,  $r_i$  and  $s_i$  denote the prizes of  $\mathcal{R}$  and  $\mathcal{S}$  for box  $i$ , respectively. Here,  $r_E$  denotes the expected prize of  $\mathcal{R}$ . What are the relevant boxes for  $\mathcal{S}$ ? Which of these boxes is  $\mathcal{R}$  willing to accept? What is the optimal weighting of the signals for  $\mathcal{S}$ ?*
- Assume  $\mathcal{S}$  could choose the best box  $i$  for her by herself (instead of sending a signal to  $\mathcal{R}$ ). Show that the ratio of the expected rewards of  $i$  and  $\varphi^*$  for  $\mathcal{S}$  can be up to  $n$ , and this is the worst case.

	1	2	3	4	5	6	7	8
$r_i$	1	3	6	7	12	13	14	16
$s_i$	8	1	10	4	8	5	6	2



**Exercise 7.3** *Online PERSUADE*

(4 + 1 + 4 + 4 points)

Consider  $n$  boxes and assume that all boxes and distributions  $\mathcal{D}_i$ ,  $i = 1, \dots, n$ , are known. The boxes are opened in a known order. In round  $i$ ,  $\mathcal{S}$  opens box  $i$  and sends a signal to  $\mathcal{R}$  (“take box  $i$ ” or “do not take box  $i$ ”). If  $\mathcal{R}$  does not take the box, the process continues with the next box. Otherwise, if  $\mathcal{R}$  takes the box, the process stops. Note that when sending the signal in round  $i$ ,  $\mathcal{S}$  only knows the content of boxes  $1, \dots, i$ .

- a) For IID boxes, show that there exists a direct and persuasive  $(1 - 1/e)^{-1}$ -competitive algorithm.  
*Hint: Use a modification of Algorithm 12 from the lecture notes where, instead of a random box, the modified algorithm takes either the first yes-box or the last no-box (the rest remains unchanged).*
- b) For independent (but not necessarily identical) boxes and the SSQ condition fulfilled with the SSQ box in round  $i^* = n$ , show that there exists an algorithm with constant competitive ratio.
- c) For independent (but not necessarily identical) boxes and the SSQ condition fulfilled with the SSQ box in round  $i^* < n$ , show that there is no algorithm with finite competitive ratio.
- d) There exists an algorithm  $\Gamma$  which computes the optimal signaling scheme in polynomial time for the independent case (even if the SSQ condition does not hold). It operates via backwards induction and solves  $n - 1$  linear programs. More precisely,  $\Gamma$  uses a LP to compute the optimal mechanism in each round  $i = 1, \dots, n$ .

Depict the LP for round  $i$ . Show that it always has a feasible solution.

*Hint: Suppose we have computed the optimal mechanism to be applied in rounds  $i + 1, \dots, n$ . Now consider round  $i$ . Under what circumstances does  $\mathcal{S}$  send “take box  $i$ ” to  $\mathcal{R}$ ? In what cases would  $\mathcal{R}$  take the recommended box indeed? Bear in mind that the mechanism signals at most one “take it”-signal in total since the scheme is supposed to be persuasive.*

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The assignments and further information on the course are provided on our website:  
<http://algo.cs.uni-frankfurt.de/lehre/oau/sommer21/oau21.shtml>

Contacts for submissions and questions: {koglin,wilhelmi}@em.uni-frankfurt.de.